AREA-TIME EFFICIENT SCALING- FREE CORDIC USING GENERALIZED MICRO-ROTATION SELECTION

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*Abstract***— This paper presents an area-time efficient CORDIC algorithm that completely eliminates the scale-factor. By suitable selection of the order of approximation of Taylor series the proposed CORDIC circuit meets the accuracy requirement, and attains the desired range of convergence. Besides we have proposed an algorithm to redefine the elementary angles for reducing the number of CORDIC iterations. A generalized micro-rotation selection technique based on high speed most significant-1-detection obviates the complex search algorithms for identifying the micro-rotations. The proposed CORDIC processor provides the flexibility to manipulate the number of iterations depending on the accuracy, area and latency requirements. Compared to the existing recursive architectures the proposed one has 17% lower slice-delay product on Xilinx Spartan XC2S200E device.**

Keywords— **coordinate rotation digital computer (CORDIC), cosine/ sine, field-programmable gate array (FPGA), most significant-1, recursive architecture.**

I. INTRODUCTION

The coordinate rotation digital computer (CORDIC) has established its popularity in several important areas of application, like generation of sine and cosine functions, calculation of discrete sinusoidal transforms like fast Fourier transform (FFT), discrete sine/cosine transforms (DST/DCT), householder transform (HT), etc. Many variations have been suggested for efficient implementation of CORDIC with less number of iterations over the conventional CORDIC algorithm. The number of CORDIC iterations are optimized in by greedy search at the cost of additional area and time for the implementation of variable scale-factor. In efficient scalefactor compensation techniques are proposed, which adversely affect the latency/throughput of computation. Two area-time efficient CORDIC architectures have been
suggested in, which involve constant scale-factor suggested in, which involve constant multiplication for adequate range of convergence (RoC). The virtually scale-free CORDIC in also requires multiplication by constant scale-factor and relatively more area to achieve

respectable RoC. The enhanced scale-free CORDIC in combines few conventional CORDIC iterations with scalingfree CORDIC iterations for an efficient pipelined CORDIC implementation with improved RoC. However, if used for recursive CORDIC architecture, combining two different types of CORDIC iterations degrades performance. In this paper, we propose a novel scaling-free CORDIC algorithm for area-time efficient implementation of CORDIC with adequate RoC. The proposed recursive architecture has comparable or less area complexity with other existing scaling-free CORDIC algorithms. Moreover, no scale-factor multiplications are required for extending the RoC to entire coordinate space, as required.

II BRIEF OVERVIEW OF CORDIC ALGORITHM

The CORDIC algorithm operates either in, rotation mode or vectoring mode, following linear, circular or hyperbolic coordinate trajectories. In this paper, we focus on rotation mode CORDIC using circular trajectory.

A. Conventional CORDIC Algorithm

In conventional CORDIC to obtain the rotated vector, the angle of rotation " μ " is decomposed into a sequence of fixed predefined elementary rotations with variable direction. The conventional rotation mode CORDIC estimates the $(i+1)$ th intermediate rotated vector from the ith vector using circular trajectory as

$$
\begin{bmatrix} x_{i+1} \\ y_{i+1} \end{bmatrix} = \mathbf{K}_i \cdot \begin{bmatrix} 1 & -\sim \tan \mathsf{r}_i \\ \sim \tan \mathsf{r}_i \end{bmatrix} \cdot \begin{bmatrix} x_i \\ y_i \end{bmatrix}
$$

Where $\mathbf{K}_i = \cos_{i}$

The sign sequence $\mu_i \in \{1,-1\}$ is so selected that

$$
{\mathbf{w}} = \sum\nolimits{i=0}^{6} \mathbf{v} \cdot \mathbf{\Gamma} \, i
$$

Where, b is the word-length in bits.

The range of convergence of this algorithm is limited to $[-99.99^\circ, 99.99^\circ]$, which can be extended to entire

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coordinate space using the properties of sine and cosine functions, using an extra iteration for full-range rotation.

The overall scaling-factor of "b" CORDIC iterations. The scaling-factor K, after sufficiently large number of iterations converges to a constant value $K =$ 0.60725.

$$
K = \prod_{i=0}^{b} Ki = \prod_{i=0}^{b} \frac{1}{\sqrt{1 + 2^{-2i}}}
$$

B. Review of Existing Scaling Free CORDIC Algorithm

Scaling-free CORDIC was the first attempt to completely dispose of the scale-factor. Here, the sine and cosine functions were approximated to

 \sin $i = (2^{-i})$ \cos $i = 1 - 2^{-(2i+1)}$

However, the approximation imposes a restriction on the basic-shift I = $[(b - 2.585)/3]$. For 16-bit data, the basic $shift = 4$ results in extremely low range of convergence. However, modified virtually adaptive scaling-free algorithm, extends the range of convergence over the entire coordinate space and introduces an adaptive scale-factor.

III. PROPOSED ALGORITHM FOR SCALING FREE **CORDIC**

The proposed design is based on the following key ideas: 1) we use Taylor series expansion of sine and cosine functions to avoid scaling operation and 2) suggest a generalized sequence of micro-rotation to have adequate range of convergence (RoC) based on the chosen order of approximation of the Taylor series.

A. Taylor Series Approximation of Sine and Cosine Functions The Taylor expansions of sine and cosine of an angle "" are given by

 $\sin \frac{1}{2}$ = $(3!)^{-1}$. 3 + $(5!)^{-1}$. 5………. cos = $1 - (2!)^{-1}$. $2 + (4!)^{-1}$. 4.........

We have estimated the maximum error in the evaluation of sine and cosine functions for different order of approximations. The maximum percentage of error in sine and cosine functions for third order approximation is 0.0033% and 0.0168%, respectively, within the permissible CORDIC elementary angles range of $[0, 7/88]$. Therefore, we choose third order of approximation for Taylor's expansion of sine and cosine functions.

1) Representation of Micro-Rotations Using Taylor Series Approximation:

The impact of orders of approximation of Taylor series of sine and cosine functions on the micro-rotations to be used in CORDIC coordinate calculation. Both theoretical and simulation results are discussed to confirm the appropriate selection of the order of approximation. Using different orders of approximation of sine and cosine functions, we can have

properties of sine and cosine
\ntion for full-range rotation.
\n
$$
x_{i+1} = \left(1 - \frac{r_i^2}{2!}\right) \cdot x_i \cdot \left(r_s - \frac{r_i^3}{3!}\right) y_i
$$

\nor K, after sufficiently large
\nrges to a constant value K = $y_{i+1} = \left(1 - \frac{r_i^2}{2!} - \frac{r_i^4}{4!}\right) \cdot x_i \cdot \left(r_s - \frac{r_i^3}{3!}\right) x_i$
\n $\frac{1}{\sqrt{1+2^{-2i}}}$
\n $x_{i+1} = \left(1 - \frac{r_i^2}{2!} - \frac{r_i^4}{4!}\right) \cdot x_i \cdot \left(r_s - \frac{r_i^3}{3!}\right) x_i$
\n $\frac{1}{\sqrt{1+2^{-2i}}}$
\n $x_{i+1} = \left(1 - \frac{r_i^2}{2!} - \frac{r_i^4}{4!}\right) \cdot x_i \cdot \left(r_s - \frac{r_i^3}{3!}\right) x_i$
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\n $\frac{1}{\sqrt{1+2^{-2i}}}$
\n $x_{i+1} = \left(1 - \frac{r_i^2}{2!} - \frac{r_i^4}{4!}\right) \$

$$
x_{i+1} = \left(1 - \frac{\Gamma_i^2}{2!} - \frac{\Gamma_i^4}{4!} - \frac{\Gamma_i^6}{6!}\right).
$$

\n
$$
\left(\Gamma_s - \frac{\Gamma_i^3}{3!} - \frac{\Gamma_i^5}{5!} - \frac{\Gamma_i^7}{7!}\right) y_i
$$

\n
$$
y_{i+1} = \left(1 - \frac{\Gamma_i^2}{2!} - \frac{\Gamma_i^4}{4!} - \frac{\Gamma_i^6}{6!}\right).
$$

\n
$$
y_i
$$

\n
$$
\left(\Gamma_s - \frac{\Gamma_i^3}{3!} - \frac{\Gamma_i^5}{5!} - \frac{\Gamma_i^7}{7!}\right) x_i
$$

We have used for coordinate calculation for evaluating the best possible combination of approximation, which satisfies the accuracy and RoC requirements, with minimum possible hardware. In Fig. 1, we have plotted the error in magnitude estimated according to (with respect to the corresponding built-in functions of MATLAB). Since the errors resulting from the five combinations are of very small order, we prefer to use for coordinate calculation with minimum complexity.

Fig.1. Error in the coordinate values *2) Expressions for Micro-Rotations Using Taylor Series Approximation and Factorial Approximation:*

Taylor series expansion with third order of approximation, with desired accuracy and RoC requirement, cannot be used in the CORDIC shift-add iterations. To implement by shift-add operations, we need to approximate the factorial terms by the power of 2 values, replacing 3! by 2^3 in the we find

$$
\begin{bmatrix} x_{i+1} \\ y_{i+1} \end{bmatrix} = \begin{bmatrix} (1 - (2!)^{-1} \cdot \Gamma_i^2) & -(\Gamma_i - 2^{-3} \cdot \Gamma_i^3) \\ (\Gamma_i - 2^{-3} \cdot \Gamma_i^3) & (1 - (2!)^{-1} \cdot \Gamma_i^2) \end{bmatrix} \cdot \begin{bmatrix} x_i \\ y_i \end{bmatrix}
$$

 x
 x_i
 B_i
 t
 t
 t
 t
 t
 32
 6
 0.25
 0.01

In Fig. 1 only, we have plotted the error in magnitude using the approximated factorial values and exact factorial values after a CORDIC rotation for initial vector with coordinates $x = 1$ and $y = 1$. The maximum percentage of error in sine and cosine values for both third order of approximation and factorial approximation is 0.0004% and 0.0168%, respectively, within the permissible CORDIC elementary angles range of $[0, 7 \, 88]$.

B. Determination of the Basic-Shift for a Given Order of Approximation of Taylor Series Expansion

One can find that:

1) the order of approximation of Taylor series expansion of sine and cosine functions determines the basic shift to be used for CORDIC iterations, and

2) the basic-shift of CORDIC micro operation determines the range of convergence.

The expressions for the basic-shifts, the first elementary angle of rotation $\begin{pmatrix} 1 \end{pmatrix}$ and RoC for different orders of approximations for different word-length of of approximations for different word-length of implementations are as follows:

Basic-shift, s =
$$
\frac{b - \log_2(n+1)!}{(n+1)}
$$

A. Organization
In the

Where b is the wordlength, and

$$
Roc = n_1 x_{-1}
$$

 n_1 is the number of microtations, and

 $1 = 2^{-s}$

In the order of approximation, the basic-shift decreases, the first elementary angle of rotation increases and RoC is expanded. Very often inclusion of higher order terms does not have any impact on the accuracy for smaller wordlengths.

The basic-shift for third order of approximation, for 16-bit word-length is [2.854]. The RoC (with basic-shift for 16-bit) is large enough to be mapped to the entire coordinate space.

esired accuracy and RoC requirement, ne CORDIC shift-add iterations. To d operations, we need to approximate the power of 2 values, replacing 3! by	Order σ f	Basic- shift		First Elementary Angle (Radians)		RoC for $n_1 = 4$ (Radians)	
	Appro	16					
(Γ_i^2) - $(\Gamma_i - 2^{-3} \Gamma_i^3)$ (Γ_i^3) $(1 - (2!)^{-1} \Gamma_i^2)$ $\left[y_i\right]$	X.	Bi	$32-$ Bit	$16-$ Bit	$32-Bit$	$\frac{16}{\text{Bit}}$	$32-Bit$
		$\mathcal{D}_{\mathcal{L}}$	6	0.25	0.01562		0.0625
nly, we have plotted the error in				0.5	0.03125	$\mathcal{D}_{\mathcal{L}}$	0.125
pproximated factorial values and exact			3	0.5	0.125		0.5

TABLE III BIT REPRESENTATION OF ELEMENTARY ANGLES AND CORRESPONDING SHIFTS

IV. GENERALIZED MICRO-ROTATION SELECTION

In the proposed generalized micro-rotation sequence, we perform multiple iterations of basic-shift, followed by non-repetitive unidirectional iterations of the micro-rotations corresponding to other shift indices, to minimize the number of iterations and achieve adequate range of convergence.

$\left\lfloor \frac{b - \log_2(n+1)!}{b - \log_2(n+1)!} \right\rfloor$ *A. Organization of Micro-Rotation Sequence*

 $n+1$ \Box In the proposed scheme, we represent the rotation angle ", " as

$$
{n} = n{1}. \quad {}_{s} + \sum_{i=1}^{n2} \Gamma_{si} . n \quad n = n_{1} + n_{2}
$$

Where ϵ is the elementary angle corresponding to the basic-shift, s_i are elementary angles for other shifts, n_1 and n_2 are non-negative integers and n represents the total number of iterations. If we do not use any micro-rotation of angle ϵ then n_1 is zero, and $n_2 = n$. On the other hand, if the

desired angle of rotation " μ " is a multiple of s then n₂ is zero and $n_1 = n$.

B. Defining the Elementary Angles

The elementary angles s and si are given by $s = 2^{-s}$
and $s_i = 2^{-si}$ where, s is the basic-shift and $s_i > s$ is the shift for ith iteration. For basic-shift = 2, we can find $s_i =$ 7 /88 and for basic-shift =3, we can find $s = 7/176$. In Table II, we list the decimal and (0, 16) fixed point binary representation of the elementary angles corresponding to different shifts.

C. Generalized Micro-Rotation Sequence Identification

The micro-rotations depending on the bit representation of the desired rotation angle in radix-2 system using most-significant-1 detector. For this we restrict the maximum rotation angle to $/4$ radians as the entire coordinate space $[0, 2]$ can be mapped to the $[0, 4]$ using octant symmetry of sine and cosine functions.

TABLE IIIII PSEUDO CODE FOR GENERATING THE MICRO-ROTATION **SEQUENCE**

Input: angle to be rotated $\binom{n}{r}$	Fig. 3 .
Begin	
M=Most Significant-1 Location of $\mu_{\rm i}$	
if $(M = 15)$ then	
$= 0.25$ radians	
shift, $s_i = 2$ and $n_{i+1} = n_{i-1}$	
Else	
Shift, $s_i = 16 - M$	Социтам
$_{n-i+1} = _{n-i}$ with $_{n-i}$ [M = '0'	
End	

If the most-significant-1 location (M) , of the rotation angle " μ " is smaller than the basic-shift "s", elementary angle of the basic-shift would be used for the CORDIC iteration. For a fixed word-length of N-bit, the shift (s_i) for the elementary angle is given by $S_i = N - M$

D. Number of Iterations to Have Desired RoC

We decide on a suitable value of "n" for realizing rotations by angles in the range $[0, 4]$. The basic-shift for 16-bit word-length is "2.854". But the basic-shift should be an integer, so we design the iterations for both "2" and "3". With basic-shift = 2 ($s = 7/88$) no more than three iterations of \overline{s} are required; therefore, the maximum value of $\overline{n_1}$ is 3. The iterations corresponding to n_2 depend on the accuracy requirements. With various values of n_2 the accuracy varies and is different for "x" and "y" coordinates.

In Fig. 2, we plot the error in the "y" coordinates estimated using for the initial vector with coordinates " $x=1$ " and "y=1", for n=7 and 8. We observe that no better bit error position (BEP) is obtained for n=8 Therefore, to minimize the number of iterations, we restrict the maximum value of "n" to 7. The error in the "x" coordinate is nearly of same magnitude as that of "y" coordinate error.

Fig. 4. Recursive architecture of the proposed CORDIC processor.

Similarly, for basic-shift = 3 ($s=7/176$), no more than six iterations of α , are equired; therefore, n=10. In Fig. 3, BEP for basic-shift $= 2$ and 3 are compared.

V. PROPOSED CORDIC ARCHITECTURE

The block diagram for the proposed CORDIC architecture is shown in Fig. 4. It makes use of the same stage for all the iterations for the coordinate calculations, as well as for the generation of shift values. The structure of each stage (shown in Fig. 5) consists of three computing blocks namely: the 1) shift-value estimation; 2) coordinate calculation; and 3) micro-rotation sequence generator.

The combinatorial circuit for the evaluation of desired shift values is shown in Fig. 6; the coordinate calculation is implemented according to (6); the combinatorial circuit for generating the micro-rotation sequence is shown in Fig. 7.

Fig. 6. Combinatorial circuit for generating the shift values.

Fig. 8. Dynamic power dissipation for the proposed architecture.

TABLE IVV SLICE DELAY PRODUCT COMPARISON FOR DIFFERENT APPROCHES

Design	No. of Slices (A)	Max. Freq. MHz (B)	Worst Case Iter. (C)	Slice Delay Product $(A*C/B)$	Max Error (%)
ALGO-I [9]	186	54.35	10	34.2	0.79
ALGO- II[9]	203	60.80	10	33.4	0.79
Scale- Free $[10]$	945	52.54	15	269.85	
Proposed	231	58.37	7	27.7	0.45

The number of iterations required in a CORDIC processor decides the rollover count of the counter. The rollover count is seven for basicshift $=2$ and ten for basicshift =3. The expiry of the counter signals the completion of a CORDIC operation; depending on this signal, the multiplexer either loads a new data-set (rotation angle, initial value of "x" and "y") to start a fresh CORDIC operation, or recycles the output of the stage to begin a new iteration for the current CORDIC operation. The input and output register files act as latches for synchronization.

VI. FPGA IMPLEMENTATION

The proposed architecture is coded in Verilog and synthesized using Xilinx ISE9.2i to be implemented in Xilinx Spartan 2E (XC2S200EPQ208-6) device. Slice-delay-product of the proposed architecture is compared with the existing CORDIC designs in Table IV; where, all designs are synthesized on Xilinx Spartan 2E XC2S200E device to maintain uniformity. The power dissipation of the proposed architecture for different clock frequencies is estimated by Xilinx XPower tool, and plotted in Fig. 8.

VII. CONCLUSION

The proposed algorithm provides a scale-free solution for realizing vector-rotations using CORDIC. The order of Taylor series approximation is decided appropriately by the proposed algorithm, not only to meet the accuracy requirement but also to attain adequate range of convergence. The generalized micro-rotation selection technique is suggested to reduce the number of iterations for low latency implementation. Moreover, a high speed most-significant-1 detection scheme obviates the complex search algorithms for identifying the micro-rotations. The proposed CORDIC processor has 17% lower slice-delay product with a penalty of about 13% increased slice consumption on Xilinx Spartan 2E device.

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